Suppose that we want produce a pulse on a digital line that was exactly 500 ms in length?

What would the code look like?

```
// Assume it is pin 0 of port B
PORTB = PORTB | 1;
delay_ms(500);
PORTB = PORTB & ~1;
```

```
// Assume it is pin 0 of port B
PORTB = PORTB | 1;
delay_ms(500);
PORTB = PORTB & ~1;
```

This will work, but why is it undesirable?

This will work, but why is it undesirable?

delay_ms() is implemented by using a
for() loop

- The microcontroller can't do anything else while it is looping
- Have to loop a precise number of times (not always easy to do)

Timing of Events: Another Example

Suppose we would want to measure the width of a pulse. How would we implement this?

Timing of Events: Another Example

```
How would we implement this?
// Wait for pin to go high
while (PINB & 0x1 == 0) { };
// Now count until it goes low
for (counter = 0; PINB & 0x1; ++counter)
 delay_ms(1);
// Now: counter is the width of
// of the pulse in ms
               Andrew H. Fagg: Embedded Real-
```

Time Systems: Timers/Counters

Timing of Events: Another Example

Again: the program cannot be doing anything else while it is waiting

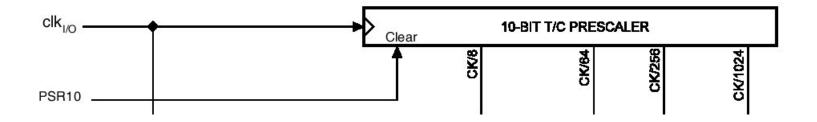
Counter/Timers in the Mega8

The mega8 incorporates three counter/timer devices in hardware.

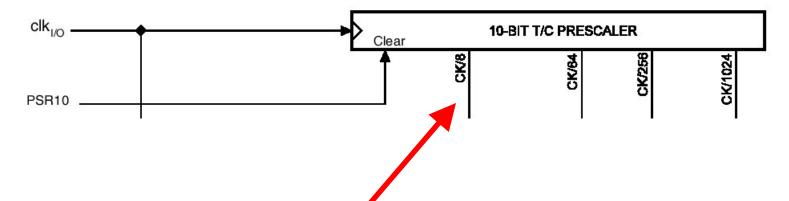
These can:

- Be used to count the number of events that have occurred (either external or internal)
- Act as a clock

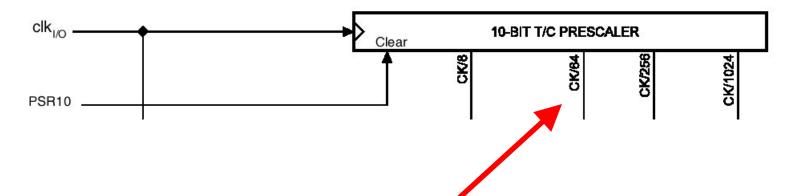
- Possible input sources:
 - Pin T0 (PD4)
 - System clock
 - Potentially divided by a "prescaler"
- 8-bit counter
- When the counter turns over from 0xFF to 0x0, an interrupt (an event) can be generated (more on this next time)



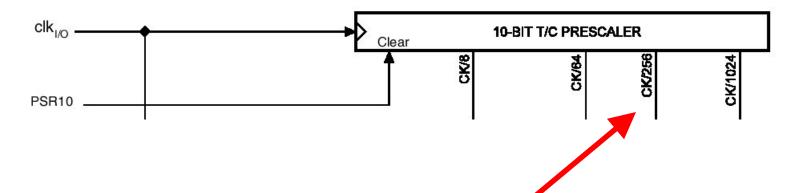
- Clock input to 10-bit counter
- Output bits: 3, 6, 8, and 10 (counting from 1)



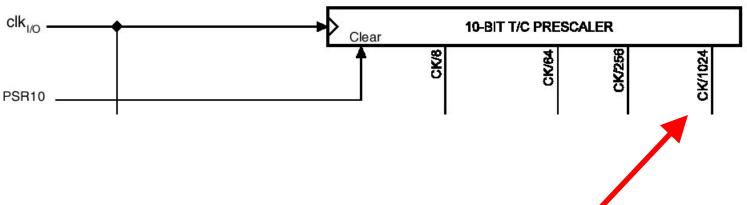
- Clock input to 10-bit counter
- Output bits: 3, 6, 8, and 10 (counting from 1)



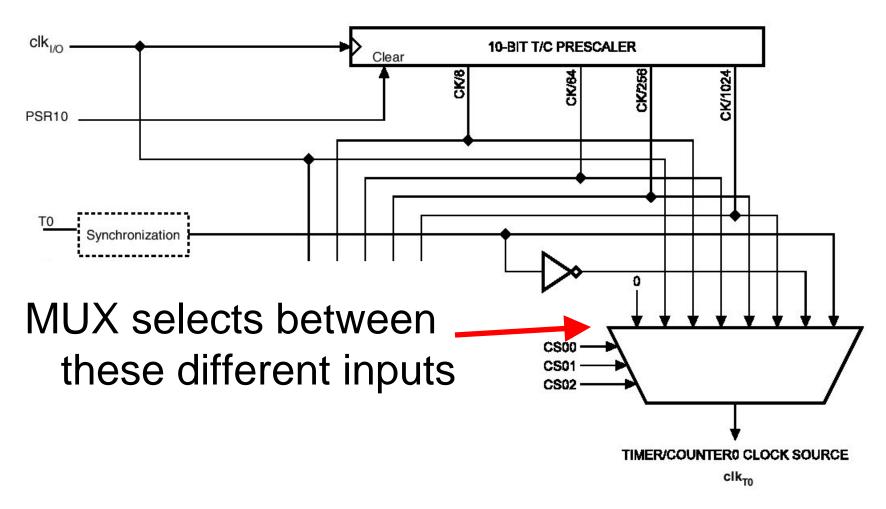
- Clock input to 10-bit counter
 Output bits: 3, 6, 8, and 10

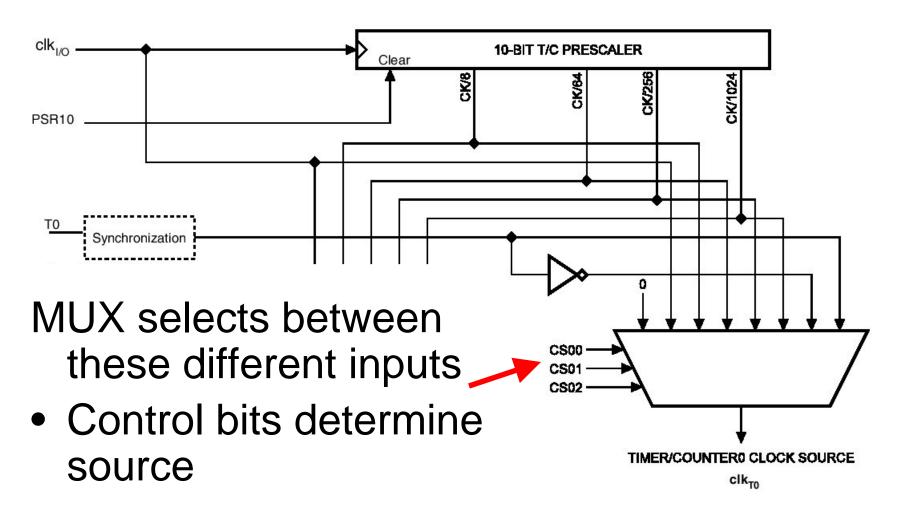


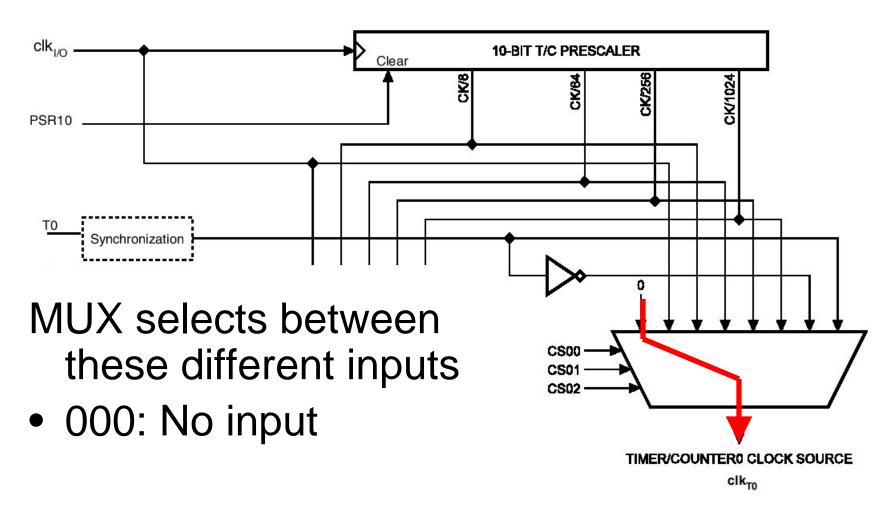
- Clock input to 10-bit counter
- Output bits: 3, 6, 8, and 10

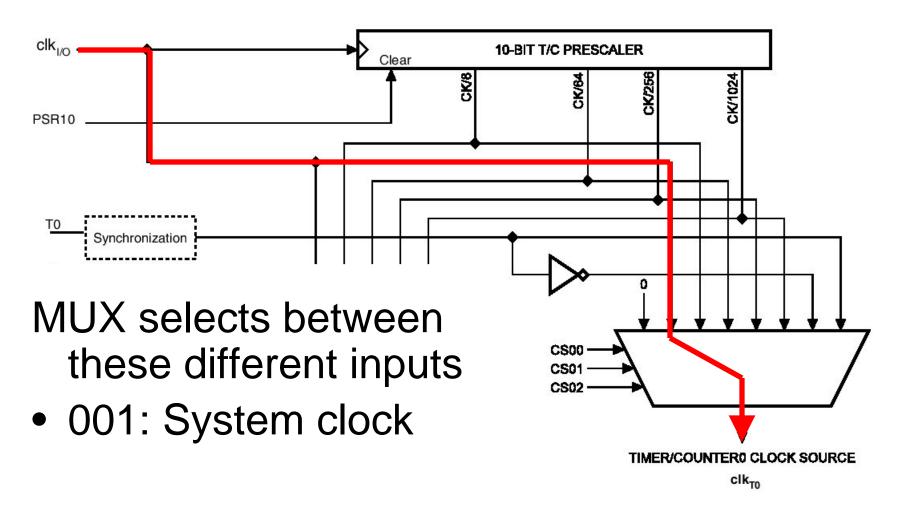


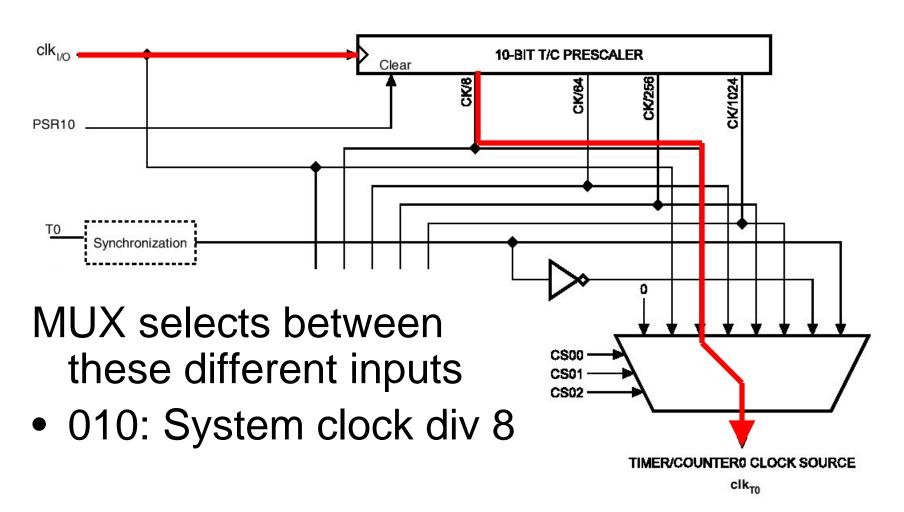
- Clock input to 10-bit counter
- Output bits: 3, 6, 8, and 10
 - These serve to divide the clock by the specified number of counts

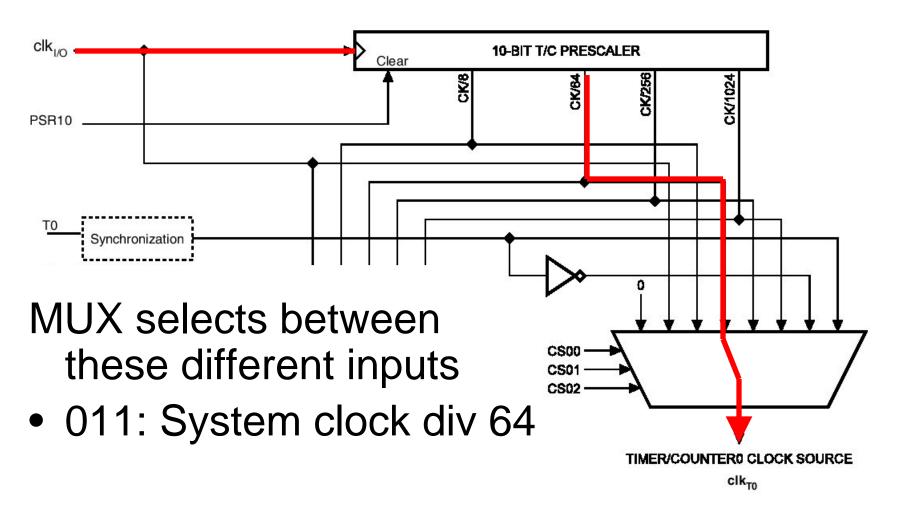


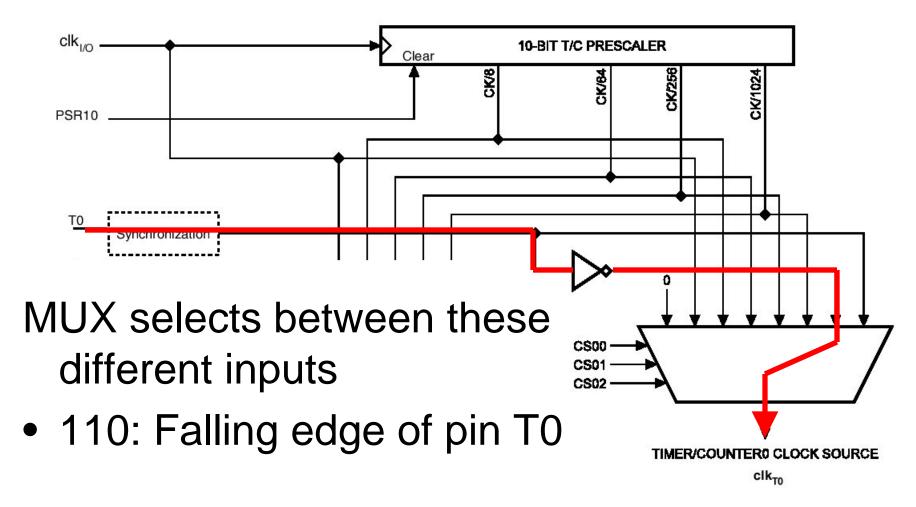


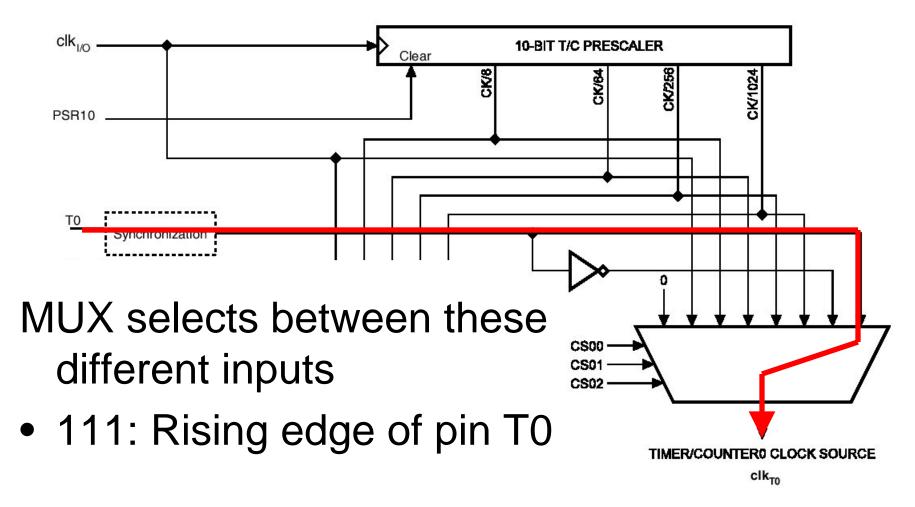


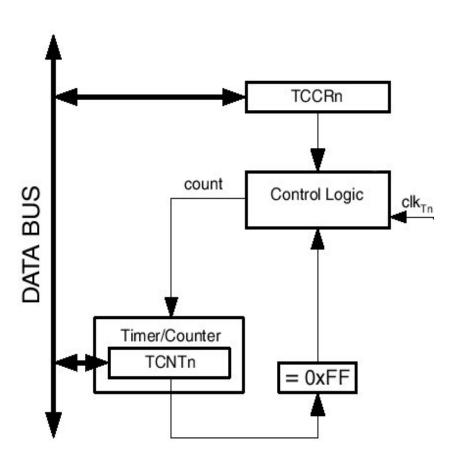




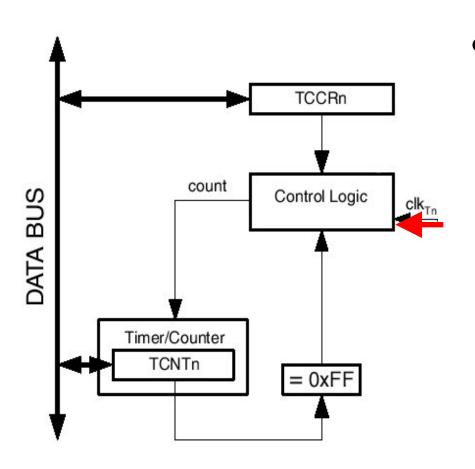




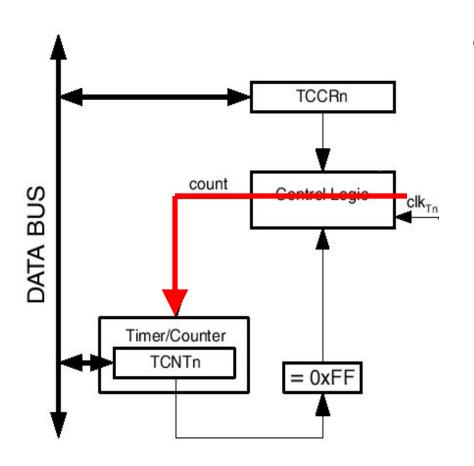




- TCNT0: 8-bit counter (a register)
- TCCR0: control register



 Clock source from previous slide



 Increment counter on every low-to-high transition

Today

- Project 3: Due April 6th
 - Pick up fan connectors
 - One group: problem compass (?)
- DC Motor control
- Timers and Interrupts
 - Timer/counter notes will be available tonight
- Tuesday: plan to spend time working on project

Timer 0 Example

Suppose:

- 16MHz clock
- Prescaler of 1024
- We wait for the timer to count from 0 to 156

How long does this take?

Timer 0 Example

$$delay = \frac{1024*156}{16,000,000} = 9948 \ \mu s \approx 10 \ ms$$

Timer 0 Code Example

```
timer0_config(TIMER0_PRE_1024); // Init: Prescale by 1024
timer0_set(0); // Set the timer to 0
<Do something else for a while>
while(timer0_read() < 156) {
  <Do something while waiting>
};
// Break out of while loop after ~10 ms
```

See Atmel HOWTO for example code (timer_demo2.c)

Timer 0 Example

Advantage over delay_ms():

- Can do other things while waiting
- Timing is much more precise
 - We no longer rely on a specific number of instructions to be executed

Timer 0 Example

One caution:

 "something else" cannot take very much time

(we have a solution for this – coming soon!)

Next Example

How do we time a delay of 100 usecs?

Next Example

How do we time a delay of 100 usecs?

```
counts*prescale = .0001*clock \_freq
= .0001*16000000
= 1600
```

Next Example

How do we time a delay of 100 usecs?

$$counts*prescale = .0001*clock _freq$$

$$= .0001*160000000$$

$$= 1600$$

$$200 * 8 = 1600$$

$$OR$$

$$25 * 64 = 1600$$

Timer 0 Code Example

```
timer0_config(TIMER0_PRE_8); // Init: Prescale by 1024
timer0_set(0); // Set the timer to 0
<Do something else for a while>
while(timer0_read() < 200) {
  <Do something while waiting>
};
// Break out of while loop after ~100 us
```

• Skip to interrupts...

Example 3: Timing the Width of a Pulse

- Input: port B, pin 1
- How long is the pin high?

timer0_config(TIMER0_PRE_1024); // Init: Prescale by 1024

```
// Wait for pin to go high
while(PINB & 0x2 == 0){};
timer0_set(0);  // Set the timer to 0
while((PINB & 0x2) != 0) {
     <Do something while waiting>
};
pulse_width = read_timer0();
```

What is the "resolution" of pulse_width?

What is the "resolution" of pulse_width?

Each "tick" of pulse_width is:

$$delay = \frac{1024}{16,000,000} = 64 \ \mu s$$

So, with pulse_width ticks:

$$delay = \frac{1024 * pulse_width}{16,000,000} = 64 * pulse_width \ \mu s$$

timer0_config(TIMER0_PRE_1024); // Init: Prescale by 1024

```
// Wait for pin to go high
while(PINB & 0x1 == 0){};
timer0_set(0);  // Set the timer to 0
while((PINB & 0x1) != 0) {
     <Do something while waiting>
};
pulse_width = read_timer0();
```

Note: the longer "something" takes, the larger the possible error in timing

Two Other Timers Besides Timer 0

Timer 1:

- 16 bit counter
- Prescalers: 1, 8, 64, 256, 1024

Timer 2:

- 8 bit counter
- Prescalers: 1, 8, 32, 64, 128, 256, 1024

Note

See oulib documentation for the list of possible prescalers for the timers

Interrupts

- Hardware mechanism that allows some event to temporarily interrupt an ongoing task
- The processor then executes a small piece of code called: interrupt handler or interrupt service routine (ISR)
- Execution then continues with the original program

Some Sources of Interrupts (Mega8)

External:

- An input pin changes state
- The UART receives a byte on a serial input

Internal:

- A clock
- Processor reset
- The on-board analog-to-digital converter completes its conversion

Interrupt Example

Suppose we are executing code from your main program:

LDS R1 (A) ← PC

LDS R2 (B)

CP R2, R1

BRGE 3

LDS R3 (D)

ADD R3, R1

STS (D), R3

Suppose we are executing code from your main program:

LDS R1 (A)

LDS R2 (B) **← PC**

CP R2, R1

BRGE 3

LDS R3 (D)

ADD R3, R1

STS (D), R3

Suppose we are executing code from your main program:

LDS R1 (A)

LDS R2 (B)

CP R2, R1 **← PC**

BRGE 3

LDS R3 (D)

ADD R3, R1

STS (D), R3

An interrupt occurs (EXT_INT1):

```
LDS R1 (A)
```

LDS R2 (B)

CP R2, R1 **← PC**

BRGE 3

LDS R3 (D)

ADD R3, R1

STS (D), R3

Execute the interrupt handler

LDS R1 (A)

LDS R2 (B)

CP R2, R1

► BRGE 3

LDS R3 (D)

ADD R3, R1

STS (D), R3

remember this location

Execute the interrupt handler

EXT_INT1:

LDS R1 (A)

LDS R2 (B)

CP R2, R1.

► BRGE 3 **¬**

LDS R3 (D)

ADD R3, R1

STS (D), R3

PC LDS R1 (G)

LDS R5 (L)

ADD R1, R2

-

RETI

Execute the interrupt handler

EXT_INT1:

LDS R1 (A)

LDS R2 (B)

CP R2, R1

► BRGE 3

LDS R3 (D)

ADD R3, R1

STS (D), R3

LDS R1 (G)

PC → LDS R5 (L)

ADD R1, R2

-

RETI

Execute the interrupt handler

EXT_INT1:

LDS R1 (A)

LDS R2 (B)

CP R2, R1

► BRGE 3

LDS R3 (D)

ADD R3, R1

STS (D), R3

LDS R1 (G)

LDS R5 (L)

PC -ADD R1, R2

•

RETI

Execute the interrupt handler

EXT_INT1:

LDS R1 (A)

LDS R2 (B)

CP R2, R1

► BRGE 3

LDS R3 (D)

ADD R3, R1

STS (D), R3

LDS R1 (G)

LDS R5 (L)

ADD R1, R2

PC -

RETI

Return from interrupt

EXT_INT1:

LDS R1 (A)

LDS R2 (B)

CP R2, R1

► BRGE 3

LDS R3 (D)

ADD R3, R1

STS (D), R3

LDS R1 (G)

LDS R5 (L)

ADD R1, R2

•

PC → RETI

Return from interrupt

LDS R1 (A)

LDS R2 (B)

CP R2, R1

►BRGE 3 ← PC

LDS R3 (D)

ADD R3, R1

STS (D), R3

EXT_INT1:

LDS R1 (G)

LDS R5 (L)

ADD R1, R2

-

RETI

Continue execution with original

EXT_INT1:

LDS R1 (A)

LDS R2 (B)

CP R2, R1

BRGE 3

LDS R3 (D) ← PC

ADD R3, R1

STS (D), R3

LDS R1 (G)

LDS R5 (L)

ADD R1, R2

•

RETI

Continue execution with original

EXT_INT1:

LDS R1 (A)

LDS R2 (B)

CP R2, R1

BRGE 3

LDS R3 (D)

ADD R3, R1← PC

STS (D), R3

LDS R1 (G)

LDS R5 (L)

ADD R1, R2

•

RETI

Interrupt Routines

Generally a very small number of instructions

- We want a quick response so the processor can return to what it was originally doing
- No delays, waits, or floating point operations in the ISR...

Timer 0 Interrupt

We can configure the timer to generate an interrupt every time that the timer's counter "rolls over" from 0xFF to 0x00

Timer 0 Interrupt Example

Suppose:

- 16MHz clock
- Prescaler of 1024

How often is the interrupt generated?

Timer 0 Example

$$interval = \frac{1024 * 256}{16,000,000} = 16.384 \, ms$$

Timer 0 Interrupt Service Routine (ISR)

An ISR is a type of function that is called when the interrupt is generated

```
ISR(TIMER0_OVF_vect) {
    // Toggle the LED attached to bit 0 of port B
    PORTB ^= 1;
};
```

What is the flash frequency?

Timer 0 Interrupt Service Routine (ISR)

```
ISR(TIMER0_OVF_vect) {
    // Toggle the LED attached to bit 0 of port B
    PORTB ^= 1;
};
```

What is the flash frequency?

$$frequency = \frac{16,000,000}{1024*256*2} = 30.5176 Hz$$

Example I: ISR Initialization in Main Program

```
// Interrupt occurs every (1024*256)/16000000 = .016384 seconds
timer0_config(TIMER0_PRE_1024);
// Enable the timer interrupt
timer0_enable();
// Enable global interrupts
sei();
while(1) {
 // Do something else
};
```

Timer 0 with Interrupts

This solution is particularly nice:

- "something else" does not have to worry about timing at all
- PB0 state is altered asynchronously from what is happening in the main program

Next Example: Timer 0 Example II

$$interval = \frac{1024 * 256}{16,000,000} = 16.384 \ ms$$

How many interrupts do we need so that we toggle the state of PB0 every second?

Timer 0 Example II

How many interrupts do we need so that we toggle the state of PB0 every second?

$$counts = \frac{1000 \, ms}{16.384 \, ms} = 61.0352$$

We will assume 61 is close enough.

Example II: Interrupt Service Routine (ISR)

```
ISR(TIMER0_OVF_vect) {
    static uint8_t counter = 0;
    ++counter;
    if(counter == 61) {
        // Toggle output state every 61st interrupt:
        // This means: on for ~1 second and then off for ~1 sec
        PORTB ^= 1;
        counter = 0;
    };
};
```

See Atmel HOWTO for example code (timer_demo.ew H. Fagg: Embedded Real-lime Systems: Timers/Counters

Example II: Interrupt Service Routine (ISR)

```
uint8_t counter = 0;
ISR(TIMER0_OVF_vect) {
   ++counter;
   if(counter == 61) {
        // Toggle output state every 61st interrupt:
        // This means: on for ~1 second and then off for ~1 sec
        PORTB ^{1}
        counter = 0;
};
```

Example II: Initialization (same as before)

```
// Initialize counter
counter = 0;
// Interrupt occurs every (1024*256)/16000000 = .016384 seconds
timer0_config(TIMER0_PRE_1024);
// Enable the timer interrupt
timer0_enable();
// Enable global interrupts
sei();
while(1) {
  // Do something else
};
```

Timer 0 Example II

What is the flash frequency?

Timer 0 Example II

What is the flash frequency?

$$frequency = \frac{16,000,000}{1024*256*61*2} \approx 0.5 Hz$$

Skip to PWM

Interrupts and Timers

Timing can often involve a cascade of multiple counters:

- prescaler (1 ... 1024)
- Timer0 (256)
- Counter within an interrupt routine (any)

Each counter implements a frequency division

Information Encoding

Many different options for encoding information for transmission to/from other devices:

- Parallel digital
- Serial digital (Project 2)
- Analog: use voltage to encode a value

Information Encoding

An alternative: pulse-width modulation (PWM)

 Information is encoded in the time between the rising and falling edge of a pulse

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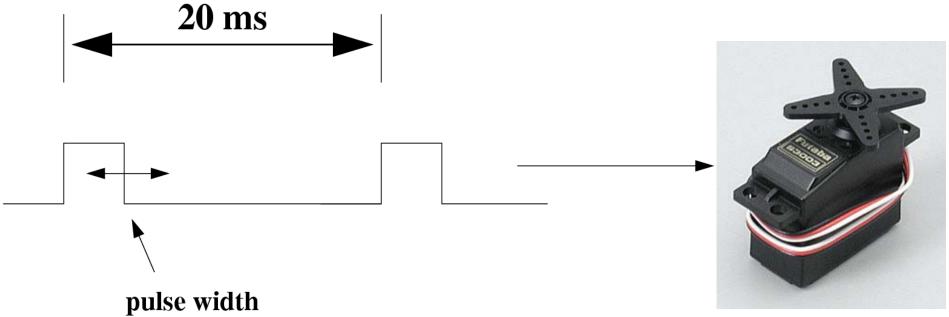
PWM Example:

RC Servo Motors

- 3 pins: power (red), ground (black), and command signal (white)
- Signal pin expects a PWM signal



PWM Example



determines motor position

Internal circuit translates pulse width into a goal position:

- 0.5 ms: 0 degrees
- 1.5 ms: 180 degrees

RC Servo Motors

- Internal potentiometer measures the current orientation of the shaft
- Uses a Position Servo Controller: the difference between current and commanded shaft position determines shaft velocity.
- Mechanical stops limit the range of motion
 - These stops can be removed for unlimited rotation

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PWM Example II: Controlling LED Brightness

What is the relationship of current flow through an LED and the rate of photon emission?

What is the relationship of current flow through an LED and the rate of photon emission?

They are linearly related (essentially)

Suppose we pulse an LED for a given period of time with a digital signal: what is the relationship between pulse width and number of photons emitted?

Suppose we pulse an LED for a given period of time with a digital signal: what is the relationship between pulse width and number of photons emitted?

- Again: they are linearly related (essentially)
- If the period is short enough, then the human eye will not be able to detect the flashes

We need:

- To produce a periodic behavior, and
- A way to specify the pulse width (or the duty cycle)

How do we implement this in code?

How do we implement this in code?

One way:

- Interrupt routine increments an 8-bit counter
- When the counter is 0, turn the LED on
- When the counter reaches some "duration", turn the LED off

Our Implementation I

```
volatile uint8_t duration = 0;
ISR(TIMER0_OVF_vect)
 static uint8_t counter = 0;
  if(duration <= counter) PORTB &= ~0x10;
  if(counter == 0) PORTB = 0x10;
 ++counter;
```

Our Implementation II

```
volatile uint8_t duration = 0;
ISR(TIMERO OVF vect)
  static uint8_t counter = 0;
  static uint8_t duration_local = 0;
  if(duration_local == counter) PORTB &= ~0x10;
  if(counter == 0) {
      duration local = duration;
      PORTB = 0x10;
  ++counter;
};
```

Another Implementation I

```
volatile uint8_t duration = 0;
ISR(TIMERO_OVF_vect)
  static uint8_t counter = 0;
  ++counter;
  if(counter >= duration)
      PORTB &= \sim 1;
  else if(counter == 0)
      PORTB |= 1;
```

Initialization Details

- Set up timer
- Enable interrupts
- Set duration in some way
 - In this case, we will slowly increase it

What does this implementation look like?

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Initialization

```
int main(void) {
  DDRB = 0x01;
  PORTB = 0;
  duration = 0;
  // Interrupt configuration
  timer0_config(TIMER0_NOPRE); // Prescaler = 1
  // Enable the timer interrupt
  timer0_enable();
  // Enable global interrupts
  sei();
```

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What is the resolution (how long is one increment of "duration")?

What is the resolution (how long is one increment of "duration")?

The timer0 counter (8 bits) expires every 256 clock cycles

$$t = \frac{256}{16000000} = 16 \ \mu s$$

(assuming a 16MHz clock)

What is the period of the pulse?

What is the period of the pulse?

The 8-bit software counter expires every 256 interrupts

$$t = \frac{256 * 256}{16000000} = 4.096 \, ms$$

Doing "Something Else"

unsigned int i; while(1) { for(i = 0; i < 256; ++i)duration = i;delay_ms(50);

ISR Example III

```
ISR(TIMER0_OVF_vect) {
  // Toggle the LED attached to bit 0 of port B
   PORTB ^{1}
};
int main(void){
   timer0_config(TIMER0_PRE_8);
   timer0_enable();
   sei();
 while(1) {
  // Do something else
                            What is the flash frequency?
 };
```

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Timer 0 Example III

What is the flash frequency?

$$frequency = \frac{16,000,000}{8*256*2} \approx 3.9 \text{ KHz}$$

ISR Example III: How about this case?

```
ISR(TIMER0_OVF_vect) {
  // Toggle the LED attached to bit 0 of port B
   PORTB ^{1}
   timer0_set(128);
};
int main(void){
   timer0_config(TIMER0_PRE_8);
   timer0_enable();
   sei();
 while(1) {
                            What is the flash frequency?
  // Do something else
 };
```

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Timer 0 Example III

What is the flash frequency?

$$frequency = \frac{16,000,000}{8*128*2} \approx 7.8 \text{ KHz}$$

Hint: key trick for project 3

3 Different Timers

- Timer 0
- Timer 1
- Timer 2

Interrupt Service Routines

- Should be very short
 - No "delays"
 - No busy waiting
 - Function calls from the ISR should be short also
 - Minimize looping
 - No "printf()"
- Communication with the main program using global variables

Interrupts, Shared Data and Compiler Optimizations

- Compilers (including ours) will often optimize code in order to minimize execution time
- These optimizations often pose no problems, but can be problematic in the face of interrupts and shared data

For example:

$$A = A + 1;$$

$$C = B * A$$

Will result in 'A' being fetched from memory once (into a general-purpose register) – even though 'A' is used twice

Now consider:

```
while(1) {
    PORTB = A;
}
```

What does the compiler do with this?

The compiler will assume that 'A' never changes.

This will result in code that looks something like this:

```
R1 = A; // Fetch value of A into register 1
while(1) {
    PORTB = R1;
}
```

The compiler only fetches A from memory once!

This optimization is generally fine – but consider the following interrupt routine:

```
ISR(TIMERO_OVF_vect) {
   A = PIND;
}
```

This optimization is generally fine – but consider the following interrupt routine:

```
ISR(TIMERO_OVF_vect) {
   A = PIND;
}
```

- The global variable 'A' is being changed!
- The compiler has no way to anticipate this

The fix: the programmer must tell the compiler that it is not allowed to assume that a memory location is not changing

 This is accomplished when we declare the global variable:

volatile uint8_t A;